

Prof. Ravi Sandhu

CONTEXT

Mid to late 90's

- SSL 1.0 never released
- SSL 2.0 flawed
- SSL 3.0 complete redesign
- > TLS from Netscape to IETF

Competitors

- SET backed by credit card companies
- S-HTTP (as opposed to https)
- > IPSEC backed by IETF committees
- SSH for secure remote access to Unix hosts

CRYPTOGRAPHIC SERVICES

Confidentiality

Encryption leaks profusely via side channels

Authentication + Integrity

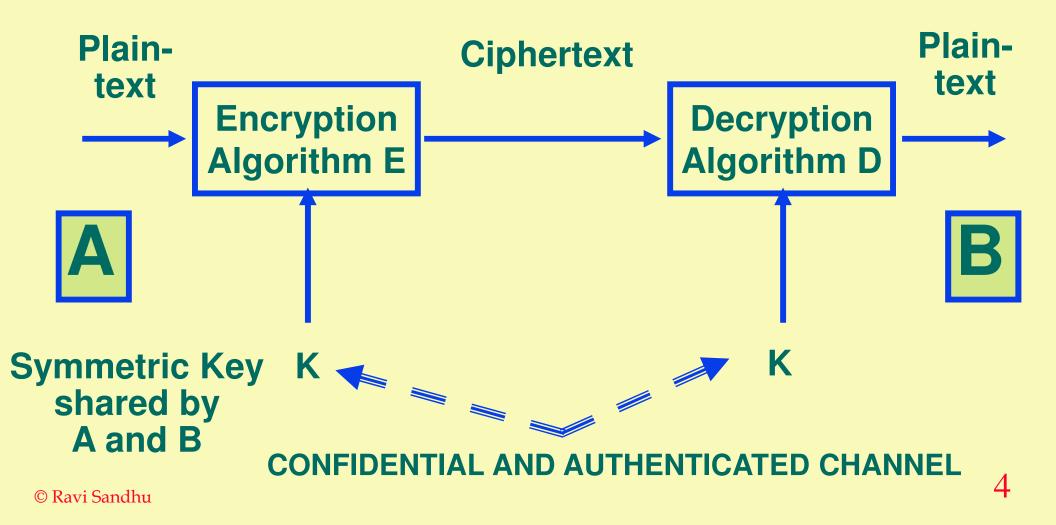
No point having one without the other

Non-repudiation

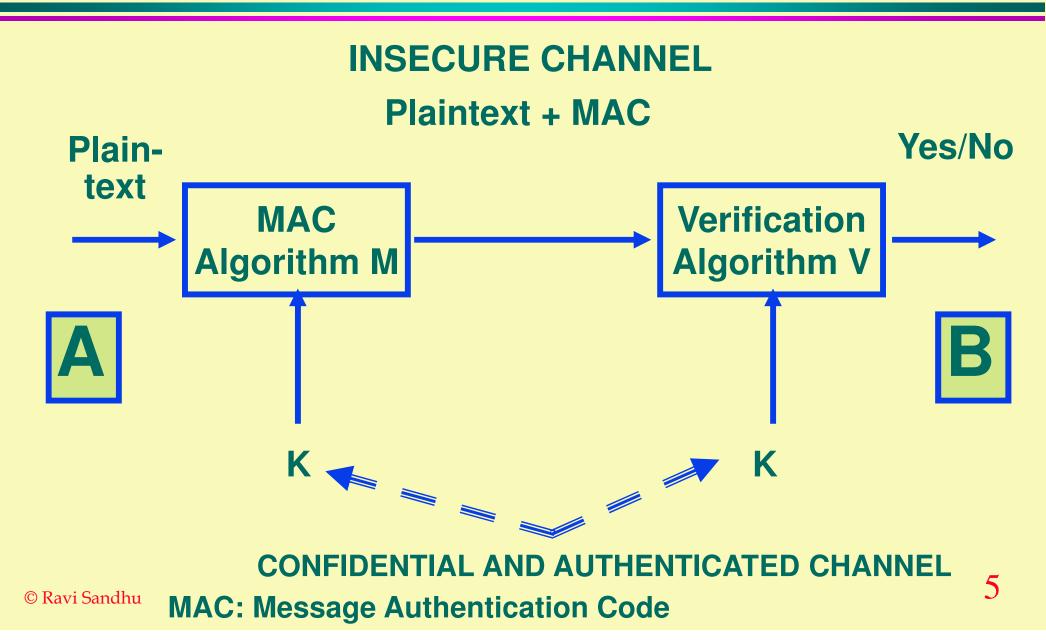
Requires asymmetric cryptography

SYMMETRIC KEY ENCRYPTION

INSECURE CHANNEL

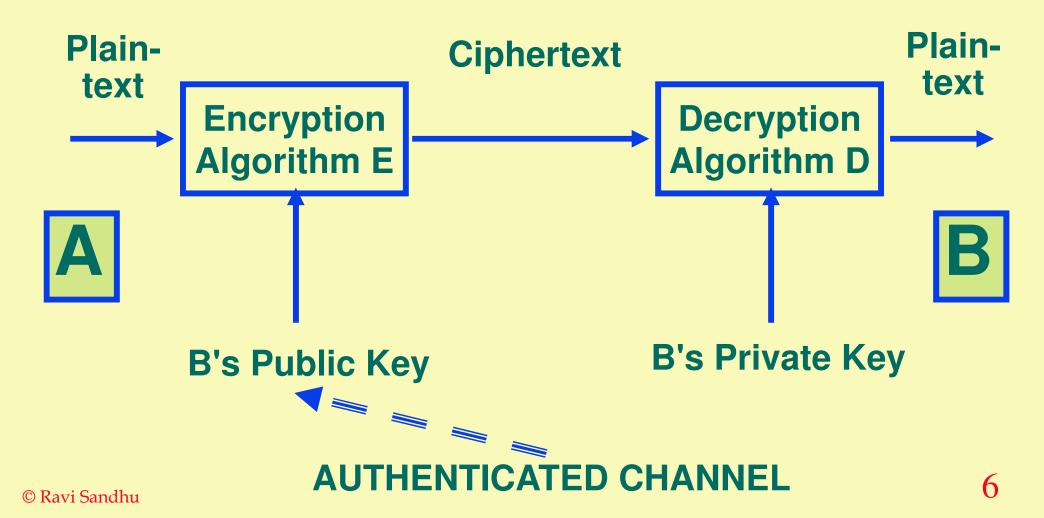


SYMMETRIC KEY AUTHENTICATION

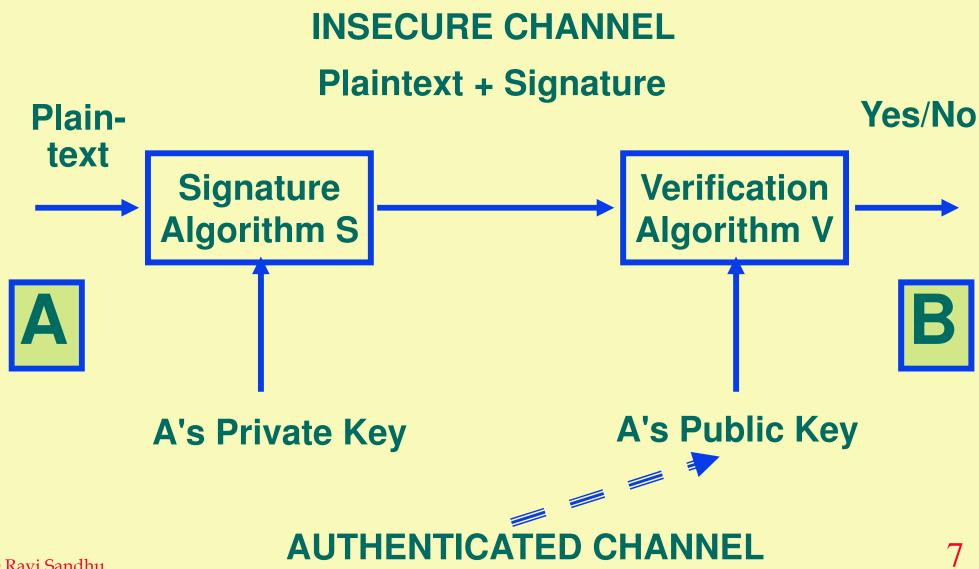


ASYMMETRIC KEY ENCRYPTION

INSECURE CHANNEL



ASYMMETRIC KEY DIGITAL SIGNATURES

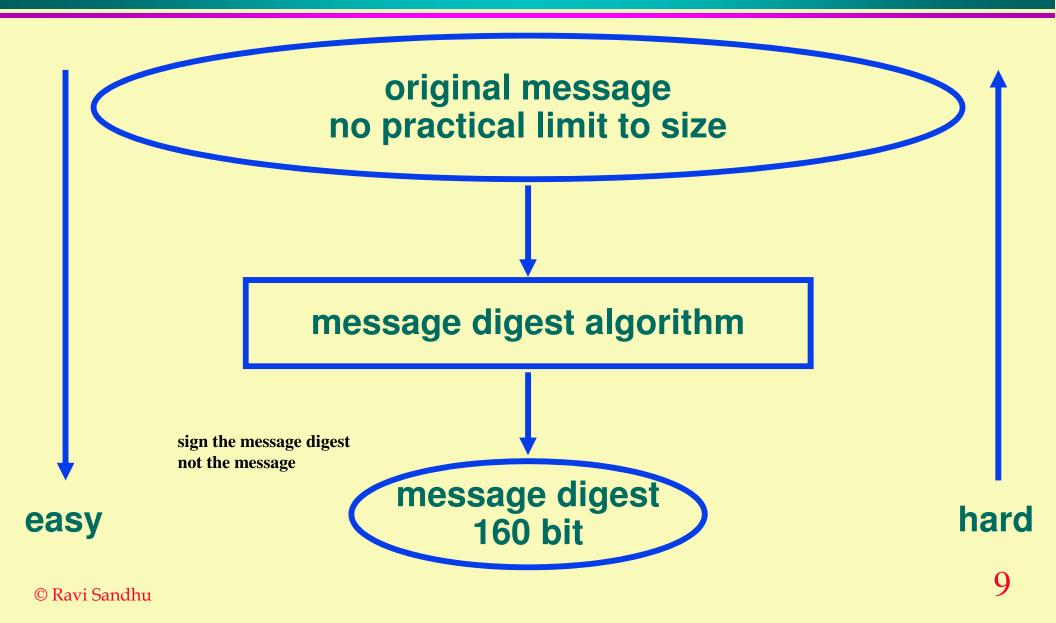


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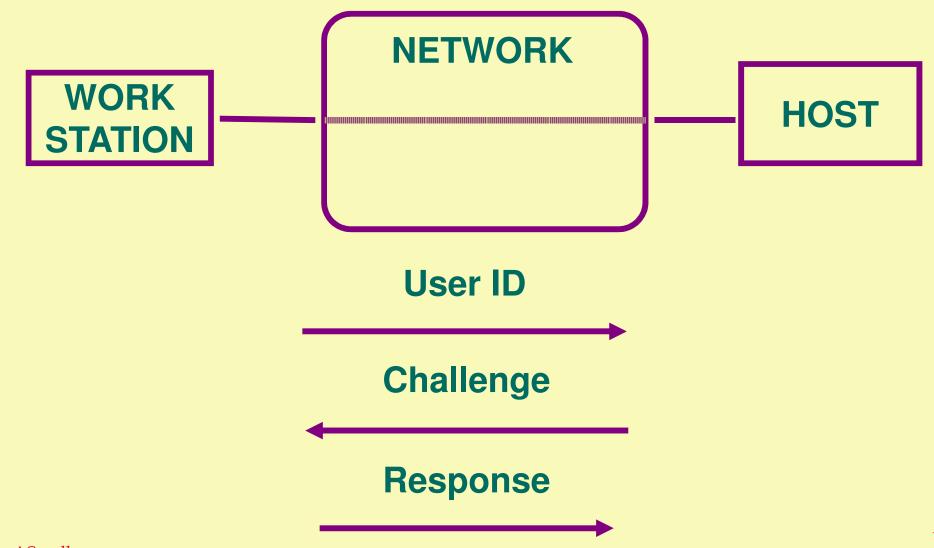
SPEED OF ASYMMETRIC KEY VERSUS SYMMETRIC KEY

- Asymmetric key runs 2-3 orders of magnitude slower than symmetric key
- This large difference in speed is likely to remain independent of technology advances

MESSAGE DIGEST



CHALLENGE RESPONSE AUTHENTICATION



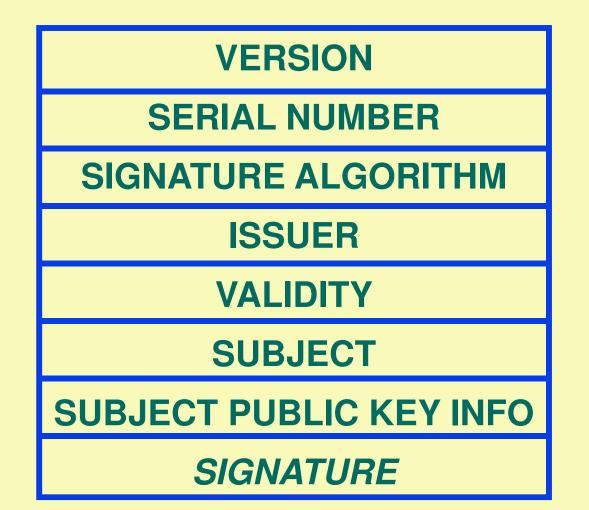
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PUBLIC-KEY CERTIFICATES

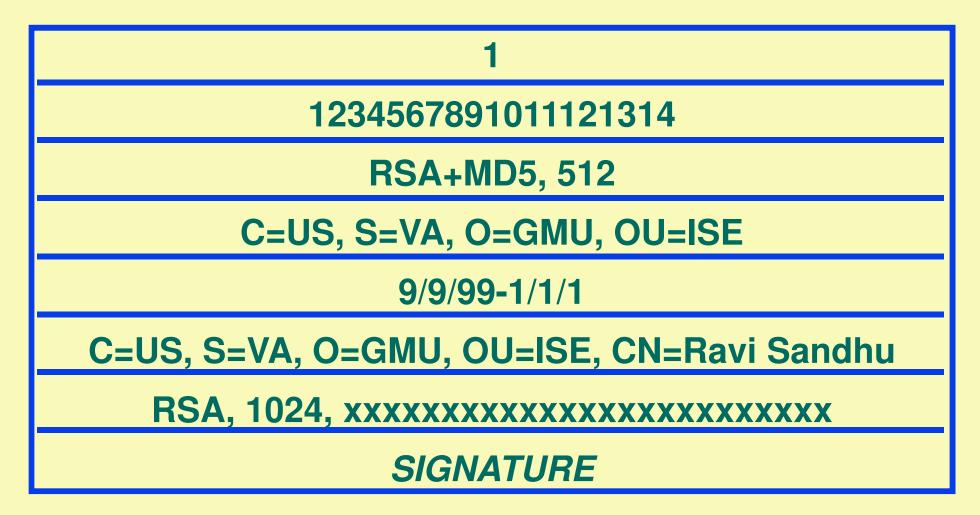
- authenticated distribution of publickeys
- * public-key encryption
 - > sender needs public key of receiver
- public-key digital signatures
 - receiver needs public key of sender
- public-key key agreement

> either one or both need the other's public key

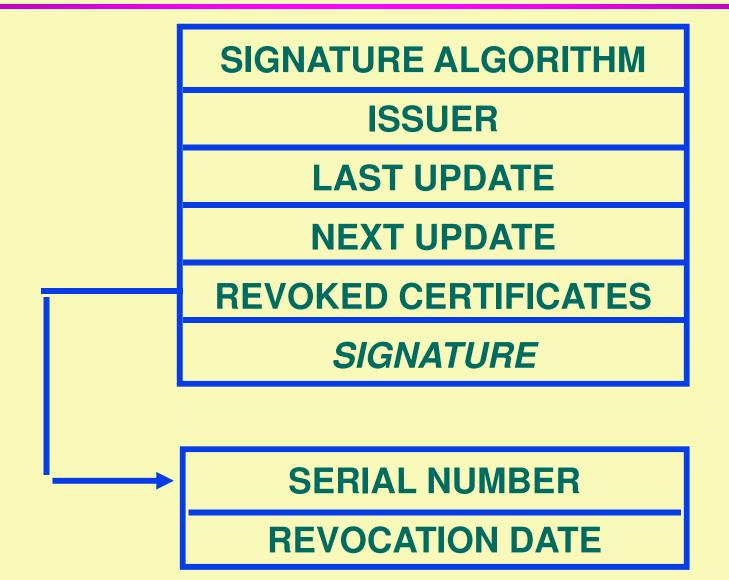
X.509v1 CERTIFICATE authenticated distribution of public-keys



X.509v1 CERTIFICATE



CRL FORMAT



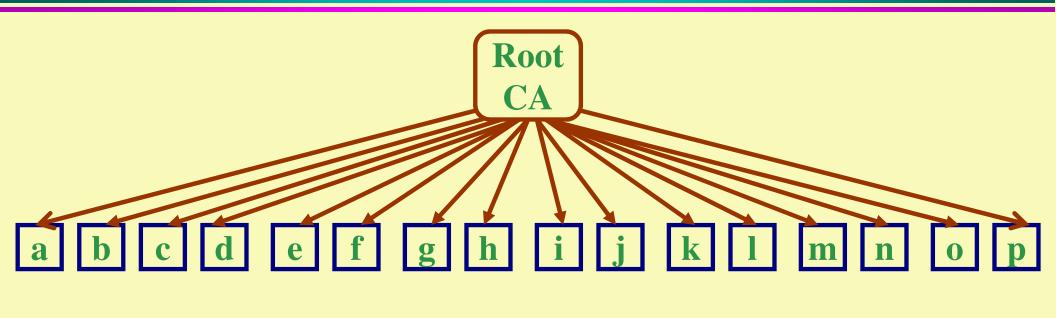
X.509 CERTIFICATES

X.509v1 very basic *** X.509v2** > adds unique identifiers to prevent against reuse of X.500 names ♦ X.509v3 > adds many extensions can be further extended

CERTIFICATE TRUST

- how to acquire public key of the issuer to verify signature
- whether or not to trust certificates
 signed by the issuer for this subject

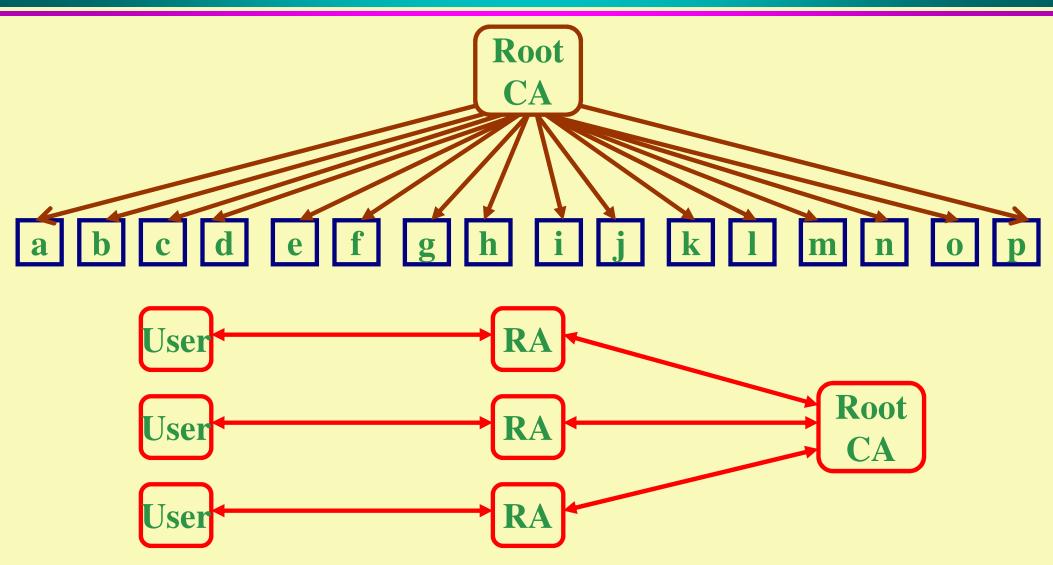
SINGLE ROOT CA MODEL



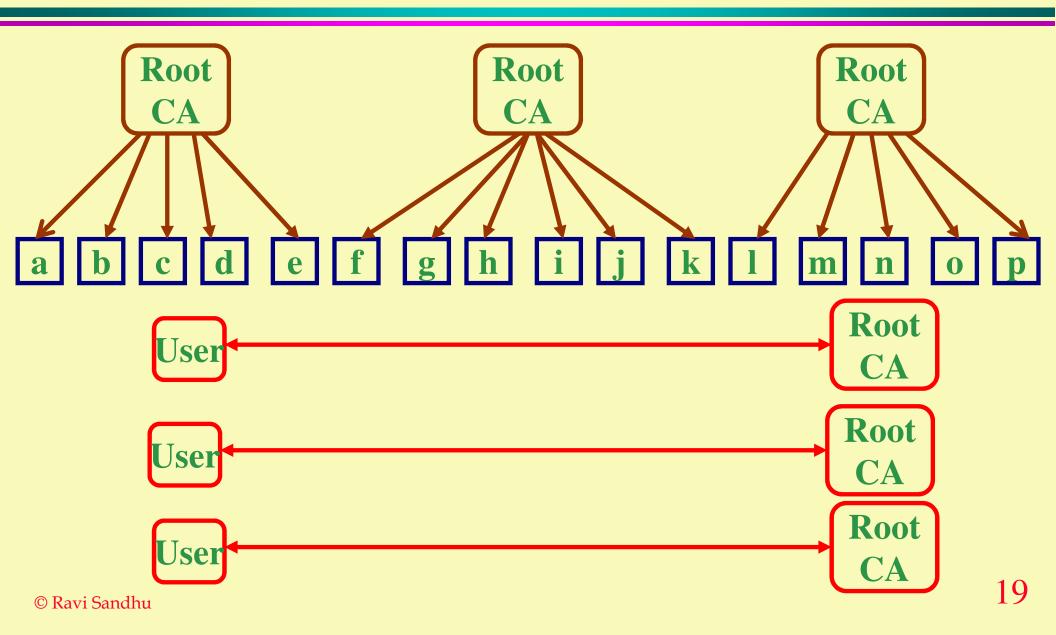


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SINGLE ROOT CA MULTIPLE RA's MODEL

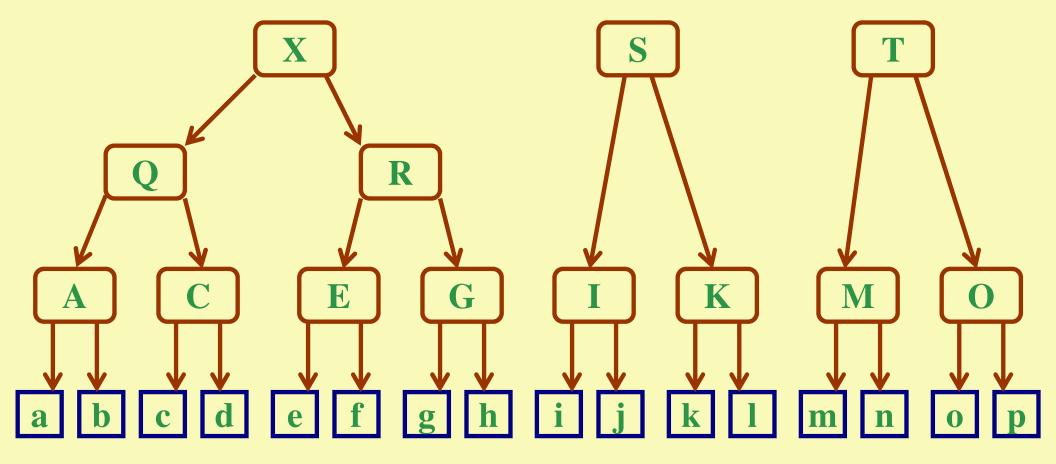


MULTIPLE ROOT CA's MODEL

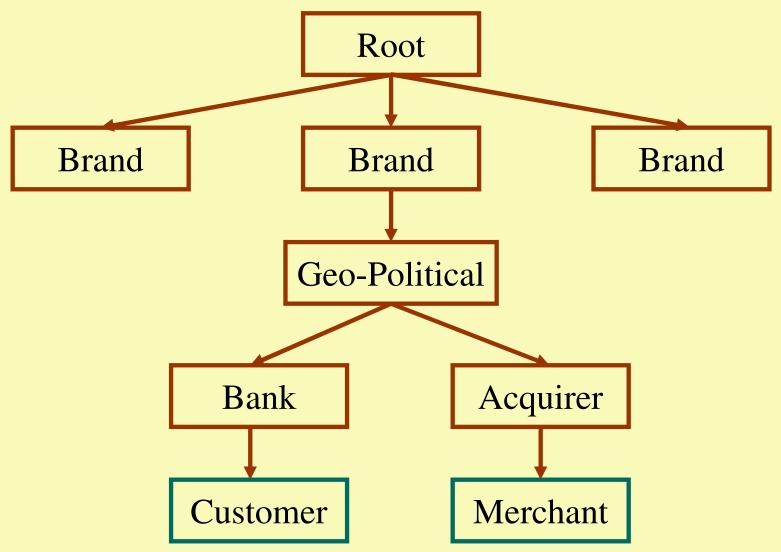


MULTIPLE ROOT CA's PLUS INTERMEDIATE CA's MODEL

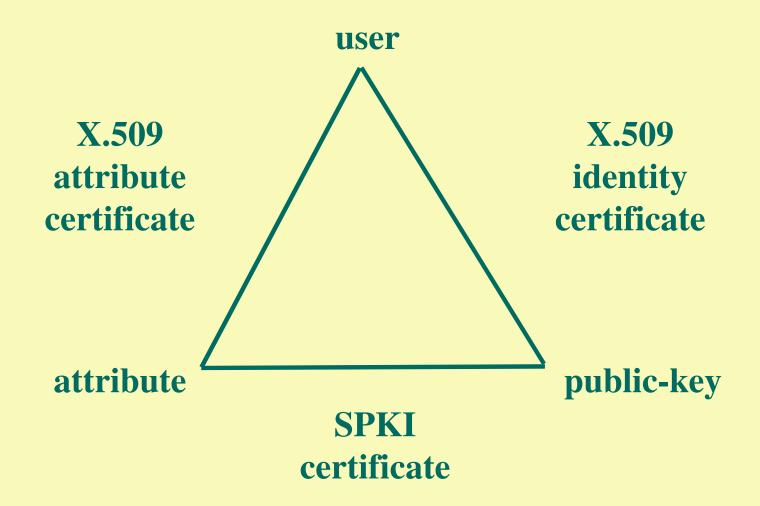
ESTABLISHED BROWSER MODEL



SECURE ELECTRONIC TRANSACTIONS (SET) CA HIERARCHY



THE CERTIFICATE TRIANGLE



SSL SERVICES

- * peer entity authentication
- * data confidentiality
- data authentication and integrity
- compression/decompression
- generation/distribution of session keys
 integrated into protocol
- security parameter negotiation

SSL ARCHITECTURE

Protocol Protocol Protocol Protocols		SSL Change Cipher Spec Protocol		HTTP	Other Application Protocols
--------------------------------------	--	---------------------------------------	--	------	-----------------------------------

SSL Record Protocol

TCP

IP

APPLICATION PORTS

https	443
ssmtp	465
snntp	563
sldap	636
spop3	995

- ftps 990
- imaps 991
- telnets 992
- ircs 993

SSL ARCHITECTURE

- Handshake protocol: complicated
 - embodies key exchange & authentication
 - > 10 message types
- Record protocol: straightforward
 - fragment, compress, MAC, encrypt
- Change Cipher Spec protocol: straightforward
 - single 1 byte message with value 1
 - could be considered part of handshake protocol
- Alert protocol: straightforward
 - > 2 byte messages
 - 1 byte alert level- fatal or warning; 1 byte alert code

SSL SESSION

SSL session negotiated by handshake protocol

- session ID
 - chosen by server
- X.509 public-key certificate of peer
 - possibly null
- compression algorithm
- cipher spec
 - encryption algorithm
 - message digest algorithm
- master secret
 - 48 byte shared secret
- is resumable flag
 - can be used to initiate new connections
 - each session is created with one connection, but additional connections within the session can be further created

SSL CONNECTION STATE

- connection end: client or server
- client and server random: 32 bytes each
- keys generated from master secret, client/server random
 - > client_write_MAC_secret server_write_MAC_secret
 - > client_write_key server_write_key
 - > client_write_IV server_write_IV
- compression state
- cipher state: initially IV, subsequently next feedback block
- sequence number: starts at 0, max 2⁶⁴-1

SSL CONNECTION STATE

4 parts to state

- current read state
- current write state
- pending read state
- pending write state
- handshake protocol
 - initially current state is empty
 - either pending state can be made current and reinitialized to empty

SSL RECORD PROTOCOL

* 4 steps by sender (reversed by receiver)

- Fragmentation
- Compression
- ► MAC
- Encryption

SSL RECORD PROTOCOL

each SSL record contains

- content type: 8 bits, only 4 defined
 - change_cipher_spec
 - alert
 - handshake
 - application_data
- protocol version number: 8 bits major, 8 bits minor
- Iength: max 16K bytes (actually 2¹⁴+2048)
- > data payload: optionally compressed and encrypted
- > message authentication code (MAC)

- initially SSL session has null compression and cipher algorithms
- both are set by the handshake
 protocol at beginning of session
- Andshake protocol may be repeated during the session

Type: 1 byte
10 message types defined
length: 3 bytes
content

Phase 1	Client		Server
	ClientHello	>	
			ServerHello
Phase 2			Certificate*
			ServerKeyExchange*
			CertificateRequest*
		<	ServerHelloDone
	Certificate*		
Phase 3	ClientKeyExchange		
	CertificateVerify*		
	[ChangeCipherSpec]		
Phase 4	Finished	>	
			[ChangeCipherSpec]
		<	Finished
	Application Data	<>	Application Data
Record	Fig. 1 - Message flo	w for a full	handshake
Protocol			
11000001	* Indicates optional or situat always sent.	ion-dependent	t messages that are not

- Phase 1:
 - Establish security capabilities
- Phase 2:
 - Server authentication and key exchange
- - Client authentication and key exchange
- Phase 4:
 - Finish

SSL 1-WAY HANDSHAKE WITH RSA

Dl 1	Client		Server	
Phase 1	ClientHello	>		
			ServerHello	
			Certificate*	
Phase 2			ServerKeyExchange*	
			CertificateRequest*	
		<	ServerHelloDone	
	-Certificate*			
Phase 3	ClientKeyExchange			
	CertificateVerify*			
	[ChangeCipherSpec]		· · · · · · · · · · · · · · · · · · ·	
Phase 4	Finished	>		
			[ChangeCipherSpec]	
		<	Finished	
	Application Data	<>	Application Data	
Record				
Recolu	Fig. 1 - Message flow for a full handshake			
Protocol				
	* Indicates optional or situation-dependent messages that are not			
	always sent.			

SSL 2-WAY HANDSHAKE WITH RSA

Phase 1	Client		Server
	ClientHello	>	
			ServerHello
			Certificate*
Phase 2			ServerKeyExchange*
Fliase Z			CertificateRequest*
		<	ServerHelloDone
	Certificate*		
Phase 3	ClientKeyExchange		
	CertificateVerify*		
	[ChangeCipherSpec]		
	Finished	>	
Phase 4			[ChangeCipherSpec]
		<	Finished
	Application Data	<>	Application Data
Record			
Recolu	Fig. 1 - Message	e flow for a full	handshake
Protocol			
	* Indicates optional or situation-dependent messages that are not		
	always sent.		

- these 9 handshake messages must occur in order shown
- optional messages can be eliminated
- 10th message explained later
 - hello_request message
- change_cipher_spec is a separate 1
 message protocol
 - Functionally it is just like a message in the handshake protocol

Client		Server
ClientHello	>	
		ServerHello
		[ChangeCipherSpec]
	<	Finished
[ChangeCipherSpec]		
Finished	>	
Application Data	<>	Application Data

Fig. 2 - Message flow for an abbreviated handshake

- hello_request (not shown) can be sent anytime from server to client to request client to start handshake protocol to renegotiate session when convenient
- can be ignored by client
 - if already negotiating a session
 - > don't want to renegotiate a session
 - client may respond with a no_renegotiation alert

Phase 1	Client		Server
	ClientHello	>	
			ServerHello
			Certificate*
Phase 2			ServerKeyExchange*
I mase 2			CertificateRequest*
		<	ServerHelloDone
	Certificate*		
Phase 3	ClientKeyExchange		
	CertificateVerify*		
	[ChangeCipherSpec]		
	Finished	>	
Phase 4			[ChangeCipherSpec]
		<	Finished
	Application Data	<>	Application Data
Record	Fig. 1 - Message fl	ow for a full	handshake
Protocol			
11000001	* Indicates optional or situa always sent.	tion-dependent	messages that are not

client hello

- > 4 byte timestamp, 28 byte random value
- session ID:
 - non-zero for new connection on existing session
 - zero for new connection on new session
- > client version: highest version
- > cipher_suite list: ordered list
- compression list: ordered list

server hello

- > 32 byte random value
- session ID:
 - new or reuse
- version
 - lower of client suggested and highest supported
- > cipher_suite list: single choice
- compression list: single choice

cipher suite

- key exchange method
 - RSA: requires receiver's public-key certificates
 - Fixed DH: requires both sides to have public-key certificates
 - Ephemeral DH: signed ephemeral keys are exchanged, need signature keys and public-key certificates on both sides
 - Anonymous DH: no authentication of DH keys, susceptible to man-in-the-middle attack
 - Fortezza: Fortezza key exchange we will ignore Fortezza from here on

cipher suite

- cipher spec
 - CipherAlgorithm: RC4, RC2, DES, 3DES, DES40, IDEA, Fortezza
 - MACAlgorithm: MD5 or SHA-1
 - CipherType: stream or block
 - IsExportable: true or false
 - HashSize: 0, 16 or 20 bytes
 - Key Material: used to generate write keys
 - IV Size: size of IV for CBC

Phase 1	Client		Server
	ClientHello	>	
			ServerHello
			Certificate*
Phase 2			ServerKeyExchange*
I mase 2			CertificateRequest*
		<	ServerHelloDone
	Certificate*		
Phase 3	ClientKeyExchange		
	CertificateVerify*		
	[ChangeCipherSpec]		
	Finished	>	
Phase 4			[ChangeCipherSpec]
		<	Finished
	Application Data	<>	Application Data
Record	Fig. 1 - Message flo	w for a full	handshake
Protocol			
11000001	* Indicates optional or situat always sent.	ion-dependent	t messages that are not

SSL HANDSHAKE: PHASE 2 SERVER AUTHENTICATION & KEY EXCHANGE

Certificate message

- server's X.509v3 certificate followed by optional chain of certificates
- required for RSA, Fixed DH, Ephemeral DH but not for Anonymous DH
- Server Key Exchange message
 - not needed for RSA, Fixed DH
 - needed for Anonymous DH, Ephemeral DH
 - needed for RSA where server has signature-only key
 - server sends temporary RSA public encryption key to client

SERVER AUTHENTICATION & KEY EXCHANGE

Server Key Exchange message

- signed by the server
- signature is on hash of
 - ClientHello.random, ServerHello.random
 - Server Key Exchange parameters
- Certificate Request message
 - request a certificate from client
 - specifies Certificate Type and Certificate Authorities
 - certificate type specifies public-key algorithm and use
- Server Done message
 - ends phase 2, always required

Phase 1	Client		Server
	ClientHello	>	
			ServerHello
			Certificate*
Phase 2			ServerKeyExchange*
I mase 2			CertificateRequest*
		<	ServerHelloDone
	Certificate*		
Phase 3	ClientKeyExchange		
	CertificateVerify*		
	[ChangeCipherSpec]		
	Finished	>	
Phase 4			[ChangeCipherSpec]
		<	Finished
	Application Data	<>	Application Data
Record	Fig. 1 - Message flo	w for a full	handshake
Protocol			
11000001	* Indicates optional or situat always sent.	ion-dependent	t messages that are not

SSL HANDSHAKE: PHASE 3 CLIENT AUTHENTICATION & KEY EXCHANGE

Certificate message

- send if server has requested certificate and client has appropriate certificate
 - otherwise send no_certificate alert
- Client Key Exchange message
 - content depends on type of key exchange (see next slide)
- Certificate Verify message
 - can be optionally sent following a client certificate with signing capability
 - signs hash of master secret (established by key exchange) and all handshake messages so far
 - provides evidence of possessing private key corresponding to certificate

SSL HANDSHAKE: PHASE 3 CLIENT AUTHENTICATION & KEY EXCHANGE

Client Key Exchange message

- > RSA
 - client generates 48-byte pre-master secret, encrypts with server's RSA public key (from server certificate or temporary key from Server Key Exchange message)
- Ephemeral or Anonymous DH
 - client's public DH value
- Fixed DH
 - null, public key previously sent in Certificate Message

SSL HANDSHAKE: POST PHASE 3 CRYPTOGRAPHIC COMPUTATION

48 byte pre master secret > RSA

- generated by client
- sent encrypted to server
- ► DH
 - both sides compute the same value
 - each side uses its own private value and the other sides public value

SSL HANDSHAKE: POST PHASE 3 CRYPTOGRAPHIC COMPUTATION

pre_master_secret: 48 bytes

PRF is composed of a sequence and nesting of HMACs

Phase 1	Client		Server
	ClientHello	>	
			ServerHello
			Certificate*
Phase 2			ServerKeyExchange*
I mase 2			CertificateRequest*
		<	ServerHelloDone
	Certificate*		
Phase 3	ClientKeyExchange		
	CertificateVerify*		
	[ChangeCipherSpec]		
	Finished	>	
Phase 4			[ChangeCipherSpec]
		<	Finished
	Application Data	<>	Application Data
Record	Fig. 1 - Message flo	w for a full	handshake
Protocol			
11000001	* Indicates optional or situat always sent.	ion-dependent	t messages that are not

SSL HANDSHAKE: PHASE 4 FINISH

- Change Cipher Spec message
 - > not considered part of handshake protocol but in some sense is part of it
- Finished message
 - > sent under new algorithms and keys
 - content is hash of all previous messages and master secret

SSL HANDSHAKE: PHASE 4 FINISH

Change Cipher Spec message

- > 1 byte message protected by current state
- copies pending state to current state
 - sender copies write pending state to write current state
 - receiver copies read pending state to read current state
- immediately send finished message under new current state

SSL HANDSHAKE: PHASE 4 FINISH

Finished message

verify_data

PRF(master_secret, finished_label, MD5(handshake_messages)+
SHA-1(handshake_messages)) [0..11];

finished_label

For Finished messages sent by the client, the string "client finished". For Finished messages sent by the server, the string "server finished".

handshake_messages

All of the data from all handshake messages up to but not including this message. This is only data visible at the handshake layer and does not include record layer headers.

SSL ALERT PROTOCOL

2 byte alert messages

- > 1 byte level
 - fatal or warning
- > 1 byte
 - alert code

SSL ALERT MESSAGES

Warning or fatal

close_notify(0), unexpected message(10), bad_record_mac(20), decryption failed(21), record overflow(22), decompression failure(30), handshake failure(40), bad certificate(42), unsupported certificate(43), certificate revoked(44), certificate_expired(45), certificate unknown(46), illegal parameter(47), unknown ca(48), access denied(49), decode error(50), decrypt error(51), export_restriction(60), protocol version(70), insufficient_security(71), internal error(80), user_canceled(90), no renegotiation(100),

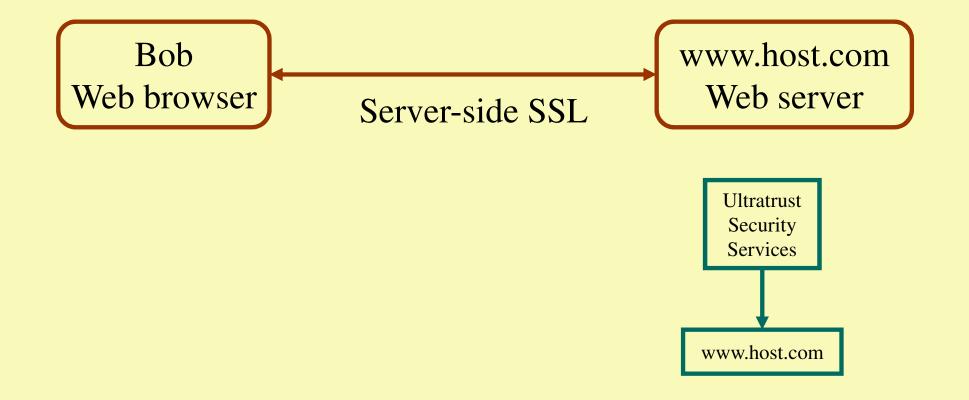
SSL ALERT MESSAGES

- always fatal
 - > unexpected_message
 - > bad_record_mac
 - > decompression_failure
 - > handshake_failure
 - > illegal_parameter

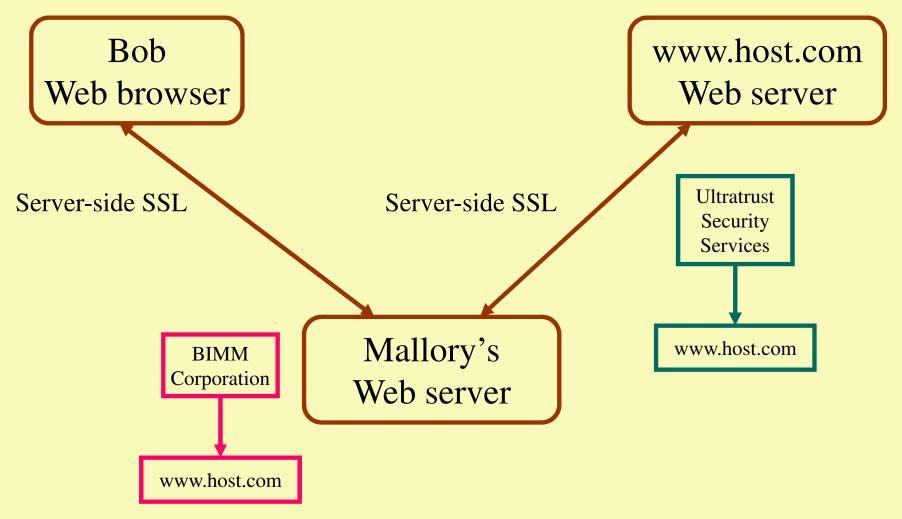
SERVER-SIDE SSL (OR 1-WAY) HANDSHAKE WITH RSA

	Client		Server
	ClientHello	>	
			ServerHello
Handshake			Certificate
Protocol		<	ServerHelloDone
	ClientKeyExchange		
	[ChangeCipherSpec] Finished	>	
	FILITSHED		[ChangeCipherSpec]
		<	Finished
	Application Data	<>	Application Data
Record			
Protocol			

SERVER-SIDE MASQUARADING

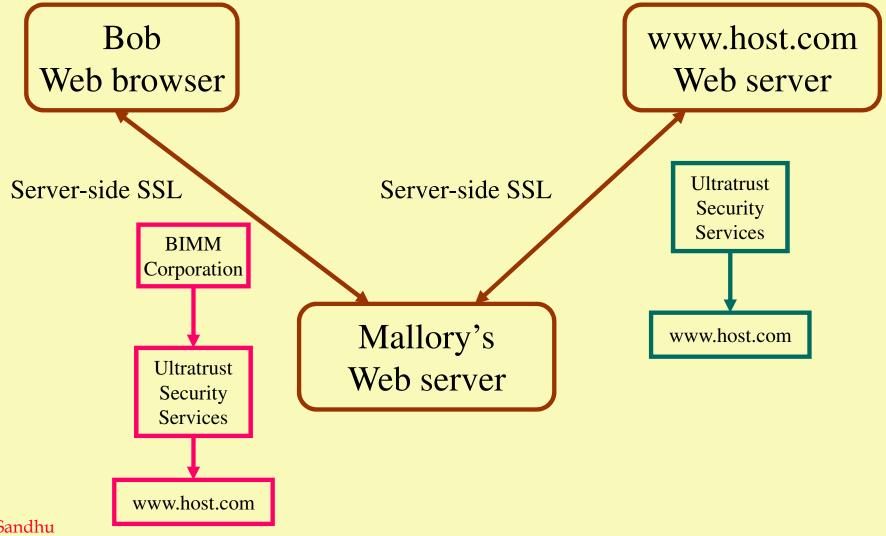


SERVER-SIDE MASQUARADING



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SERVER-SIDE MASQUARADING

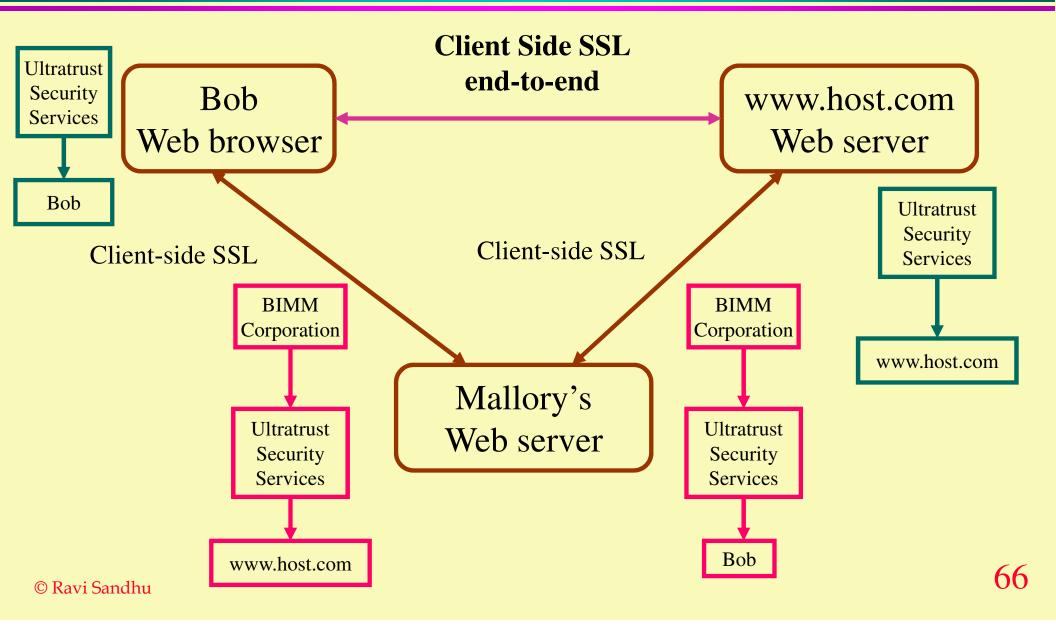


CLIENT-SIDE SSL (OR 2-WAY) HANDSHAKE WITH RSA

	Client		Server
	ClientHello	>	
			ServerHello
Handshake			Certificate
Protocol			CertificateRequest
		<	ServerHelloDone
	Certificate		
	ClientKeyExchange		
	CertificateVerify		
	[ChangeCipherSpec]		
	Finished	>	
			[ChangeCipherSpec]
		<	Finished
	Application Data	<>	Application Data
Record			
Drotocol			

Protocol

MAN IN THE MIDDLE MASQUARADING PREVENTED



SSL

- Deployed in broken form
- Guardian of e-commerce
- World's most successful crypto protocol